

# A Modified Remote User Authentication Scheme Using Smart Cards

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**Abstract**—In 2000, Hwang and Li proposed a new remote authentication scheme using smart cards. Their scheme is based on the ElGamal's public key cryptosystem. However, Chan and Cheng pointed out that the scheme is vulnerable to the masquerade attack. In this article, we shall show a different attack on Hwang-Li scheme which is easier and simpler. Furthermore, we shall present an enhanced scheme for repairing the above attacks.

**Index Terms**—Authentication, cryptography, password, security.

## I. INTRODUCTION

**P**ASSWORD authentication schemes with smart card have a history in the remote user authentication environment. A wide variety of password authentication schemes with smart card [2], [3], [7], [8], [11], [10], [12], [13], [14], [15] have been proposed. These schemes can allow a legal user to login a remote server and access the data provided by the remote server.

In 1995, Wu proposed a new remote login authentication scheme based on simple geometric Euclidean plane [16]. His scheme allows users to freely choose passwords themselves. However, Hwang [6] has pointed out that the weakness of Wu's scheme lies in the security. In 2000, Hwang and Li [9] proposed a new remote user authentication scheme using smart card based on ElGamal's cryptosystem. Chen and Chang [1] pointed out a cryptanalysis of Hwang-Li scheme's scheme.

In this article, we shall present a different attack on the Hwang-Li scheme that cannot withstand a masquerade attack; an attacker can derive a legal user's password and then masquerade as another legal user to login a remote server. For the above security flaw, we shall propose an enhanced scheme to repair.

## II. REVIEW OF THE HWANG-LI SCHEME

In this Section, we will briefly review Hwang-Li scheme [9]. The scheme is composed of three phases: the registration phase, the login phase, and the authentication phase. Each user sends his/her identity to the server in the registration phase. The server will issue a smart card and a password to every the

legal user through a secure channel when the user is identified. When the user wants to access a remote server, he/she can insert his/her smart card into the login device and then key in the identity and password to access services. The server will verify the data in the authentication phase.

- **Registration phase:** The server prepares two system parameters  $p$  and  $x_s$ , where  $p$  is a large prime number and  $x_s$  is the secret key of the server. Besides, there is also a public one-way function  $h(\cdot)$ . Suppose a new user  $U_i$  wants to register the server for accessing services. First,  $U_i$  registers his/her identity  $ID_i$  to the server. The server computes the password  $PW_i = ID_i^{x_s} \bmod p$  for the user  $U_i$  and stores  $h(\cdot)$  and  $p$  into a smart card. Then the server issues the smart card and  $PW_i$  to the user  $U_i$  through a secure channel. Note that the data stored in the smart card is the same for all users, i.e.,  $h(\cdot)$  and  $p$ . The procedure is shown in Figure 1.
- **Login phase:** The user inserts his/her smart card into the login device when he/she wants to login the server. He/she keys in his/her identity  $ID_i$  and  $PW_i$ . The smart card provides  $C_1 = ID_i^r \bmod p$  and  $C_2 = M(PW_i)^r \bmod p$  for login, where  $r$  is a random number, and  $M = ID_i^t \bmod p$ . Here,  $t = h(T \oplus PW_i) \bmod (p-1)$ , where  $T$  is the current date and time by the input device, and  $\oplus$  stands for an exclusive operation. The device sends a message  $C = (ID_i, C_1, C_2, T)$  to the remote server.
- **Authentication phase:** The server first checks  $ID_i$  and  $T$  to make sure whether they are valid. If  $(T' - T) \geq \Delta T$ , then the server rejects the login request. Here,  $T'$  is the current date and time by the server;  $\Delta T$  is the expected legal time interval for transmission delay. Second, the server checks  $C_2(C_1^{x_s}) \bmod p = (ID_i)^{h(T \oplus PW_i)}$  to see if it holds before accepting the login request.

## III. CRYPTANALYSIS OF THE HWANG-LI SCHEME

In 2000, Chan and Cheng pointed out Hwang-Li scheme cannot withstand the masquerade attack. However, we shall present an easier and simpler attack on Hwang-Li scheme. In this section, we shall briefly review Chan and Cheng's attack, and then show a different approach to attack on Hwang-Li scheme. We shall show that in the Hwang-Li scheme, it is possible for an illegal user to obtain a legal user's password and then masquerade as the legal user to login the server without being detected.

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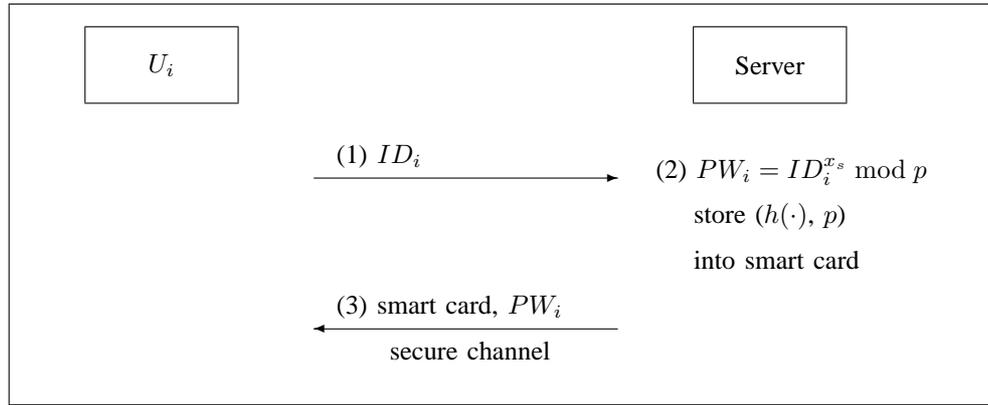


Fig. 1. The Registration Phase

### A. Chan-Cheng's attack

Suppose that a user  $U_a$  wants to masquerade other legal users. She/he submits her/his  $ID_a$  to the remote server for registering to be a legal user. The remote server will responses  $PW_a$  and a smart card for  $U_a$  after the identity is identified. Now,  $U_a$  wants to create a legal user  $U_f$  and corresponding a valid pair of  $(ID_f, PW_f)$ . Now, she/he computes  $ID_f = (ID_a \cdot ID_a) \bmod p$  and  $PW_f = ID_f^{x_s} = (PW_a \cdot PW_a) \bmod p$ . Thus,  $U_a$  can successfully login in the remote server via forged  $(ID_f, PW_f)$ .

### B. Different Attack

Suppose a user  $U_a$  is an attacker who wants to masquerade as another user  $U_i$  to login a remote server and gain access privilege.

The attacker can make his/her  $ID_a$  equal to  $ID_i^k \bmod p$  because  $ID_i$  is public, where  $k$  is a random number chosen by the attacker and  $\gcd(k, p) = 1$ . The attacker submits his/her identity  $ID_a$  and registration request to the server in the registration phase.

The server will verify the identity attached to the registration request and compute  $PW_a = ID_a^{x_s} \bmod p$  for the attacker to store  $h(\cdot)$  and  $p$  into his/her smart card. Finally, the server will issue the smart card and  $PW_a$  to the attacker through a secure channel. Now, the attacker can derive  $U_i$ 's password as follows:

$$\begin{aligned}
 (PW_a)^{-k} \bmod p &= (ID_j^{x_s})^{-k} \bmod p, \\
 &= (ID_i^{k \cdot x_s})^{-k} \bmod p, \\
 &= (ID_i^{x_s}) \bmod p, \\
 &= PW_i.
 \end{aligned}$$

Next, the attacker can successfully use  $U_i$ 's password to masquerade as the user  $U_i$  to login the remote server.

## IV. SECURITY ENHANCEMENT FOR THE HWANG-LI AUTHENTICATION SCHEME

In this section, we present a modification of the Hwang-Li scheme to enhance the security flaw described in Section 3

and analysis the security of our enhanced scheme.

### A. The Enhanced Scheme

Here, we propose a enhanced scheme that can amend the security flaw of the Hwang-Li scheme. It employs the concept of hiding identity to prevent from masquerading attack. We only modify the registration phase which issue a "shadowed" identity [4], [5] for every legal user. The steps of login and authentication phase are retained except that replace  $(ID_i, ID'_i)$  by "shadowed" identity  $(SID_i, SID'_i)$ , respectively. The modified registration phase is presented as follows.

- **Registration Phase:** Assume that this phase is executed over a secure channel. First,  $U_i$  submits her/his identity string  $J_i$  to the remote server for registration, where  $J_i$  is  $U_i$ 's identity string that includes name, unique number etc. which are unique. The remote server computes  $(SID_i, SID'_i, PW_i)$  for the registering user after her/his identity  $J_i$  is identified.

$$SID_i = Red(J_i), \quad (1)$$

$$SID'_i = SID_i^{x_{s1}} \bmod p, \quad (2)$$

$$PW_i = SID_i^{x_{s2}} \bmod p. \quad (3)$$

Where,  $Red(\cdot)$  is a "shadowed" identity of the device which only is possessed with the remote server;  $SID_i$  and  $SID'_i$  are  $U_i$ 's "shadowed" identity which can be disclosed. Furthermore, the remote server issues the smart card and  $(SID_i, PW_i)$  to  $U_i$ , which  $(f(\cdot), p)$  is stored into the smart card.

### B. The Security Analysis of the Enhanced Scheme

Our enhanced scheme is modification of the Hwang-Li scheme. The security analysis have been already discussed and demonstrated in [6]. Therefore, we shall only discuss the enhanced scheme how to resist the masquerade attack which proposed by [1] and us.

The masquerade attack on the Hwang-Li scheme is described in Section 3. The attack works because the evil user

can successfully register a new  $ID_e$  via  $ID_i$ . In our enhanced scheme, we propose a modification of the login phase as the equations (1), (2) and (3) to withstand the attack.

As the Section 3, assume that an evil user  $U_e$  can intercepts  $C = (PID_i, C_1, c_2, T)$  from a public network. Now,  $U_e$  submits her/his

$$J_e = (PID_i)^k \text{ mod } p,$$

to the remote server to register for masquerading as  $U_i$ . Upon receiving the registration message  $J_e$  from  $U_e$ , the remote server will rejects the registration request because the format of  $J_e$  is incorrect which must includes name, unique number etc. for identifying.  $J_i$  is maintained by the user  $U_i$  and the remote server, secretly. Thus,  $U_e$  cannot masquerades as  $U_i$  to login and access the remote server because cannot intercepts  $J_i$  from the public network. For this reason, our enhanced scheme can withstand the masquerading attack. As same reason, our enhanced scheme also can withstand the Chan-Cheng's attack.

## V. CONCLUSIONS

In this article, we have shown a that the Hwang-Li scheme is vulnerable to the different masquerade attack from [1]. An evil user can masquerades as another legal user to login a remote server via derive a legal user's password. For the above attacks, we shall present a enhanced scheme to repair the security flaw of the Hwang-Li scheme.

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